

## Knapsack problem

### Problem:

- input:
  - ✓ n objects.
  - ✓ each object i has a weight  $w_i$  and a profit  $p_i$
  - ✓ Knapsack : M
- output:
  - ✓ Fill up the Knapsack s.t. the total profit is maximized.
  - ✓ Feasible solution:  $(x_1, \dots, \dots, x_n)$ .

### Formally:

- ✓ Let  $x_i$  be the fraction of object i placed in the Knapsack,  $0 \leq x_i \leq 1$ .

For  $1 \leq i \leq n$ .

Then :

$$P = \sum_{1 \leq i \leq n} p_i x_i \quad \text{And} \quad \sum_{1 \leq i \leq n} w_i x_i \leq M$$

### Assumptions:

- $\sum_{i=1}^n w_i > M$  ; not all  $x_i = 1$ .
- $\sum_{1 \leq i \leq n} w_i x_i = M$

### Example:

- 3 objects ( $n=3$ ).
- $(w_1, w_2, w_3) = (18, 15, 10)$
- $(p_1, p_2, p_3) = (25, 24, 15)$
- $M=20$  (knapsack capacity)

### Largest-profit strategy: (Greedy method)

- Pick always the object with largest profit.
- If the weight of the object exceeds the remaining Knapsack capacity, take a fraction of the object to fill up the Knapsack.

**Example:**

$P=0$ (profit) ,  $C=M=20$  /\* remaining capacity \*/

- Put object 1 in the Knapsack.  
 $P=25$  Since  $w_1 < M$  then  $x_1=1$   
 $C=M-18=20-18=2$
- Pick object 2  
 Since  $C < w_2$  then  $x_2= C/w_2=2/15$ .  
 $P=25+2/15*24 =25+3.2=28.2$
- Since the Knapsack is full then  $x_3=0$ .

The feasible solution is  $(1, 2/15, 0)$   $P=28.2$

**Smallest-weight strategy:**

- be greedy in capacity: do not want to fill the knapsack quickly.
- Pick the object with the smallest weight.
- If the weight of the object exceeds the remaining knapsack capacity, take a fraction of the object.

**Example:**

$cu=M=20$

- Pick object 3  
 Since  $w_3 < cu$  then  $x_3=1$   
 $P= 15$   $cu =20-10 = 10$  ,  $x_3 =1$
- Pick object 2  
 Since  $w_2 > cu$  then  $x_2 = 10/15 = 2/3$   
 $P = 15+ 2/3*24$   
 $= 15+ 16 = 31$   $cu= 0$ .
- Since  $cu=0$  then  $x_1=0$
- Feasible solution :  $(0,2/3,1)$   $P$  (profit)=31.

**Largest profit-weight ratio strategy:**

- Order profit-weight ratios of all objects.
- $P_i/w_i \geq (P_{i+1})/(w_{i+1})$  for  $1 \leq i \leq n-1$
- Pick the object with the largest  $p/w$

- If the weight of the object exceeds the remaining knapsack capacity, take a fraction of the object.

**Example:**

$$P_1/w_1=25/18=1.389$$

$$P_2/w_2=24/15=1.6$$

$$P_3/w_3=15/10=1.5$$

$$P_2/w_2 > P_3/w_3 > P_1/w_1$$

$$C_u=20; p=0$$

- Pick object 2  
Since  $c_u \geq w_2$  then  $x_2=1$   
 $c_u=20-15=5$  and  $p=24$
- Pick object 3  
Since  $c_u < w_3$  then  $x_3=c_u/w_3=5/10=1/2$   
 $c_u=0$  and  $P=24+1/2 \cdot 15=24+7.5=31.5$
- Feasible solution  $(0,1,1/2)$   $P=31.5$

**Largest-profit strategy:** 28.2

**Smallest-weight strategy:** 31.

**Largest profit-weight ratio strategy:** 31.5

Largest profit-weight ratio strategy gives maximum profit

**Analysis**

Analysis of knapsack problem with **Largest profit-weight ratio strategy:**

**Step 1-** Calculating  $P/w$  for all  $n$  object required  $O(n)$ .

**Step 2-** Arranging  $P/w$  for all  $n$  object required sorting technique so time required for sorting  $P/w$  for all object is  $O(n \log n)$ .

**Step 3-** If the weight of the object exceeds the remaining knapsack capacity, take a fraction of the object required  $O(n)$

So **time complexity of knapsack problem** is  $O(n \log n)$  because here bigger loop exists.